Using Web N-Grams to Help Second-Language Speakers

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Writing in a foreign language is difficult.

- Problems include
 - □ Spelling
 - 🗅 Grammar
 - Translation
 - Word Choice
 - Writing Style

- Tools include
 - □ Spell checkers.
 - □ Grammar checkers.
 - Dictionaries, (machine translation).
 - Thesauri.
 - □ Style checkers.

Anything missing?

What about text commonness?

What about text commonness?

Correctness vs. Commonness

We present NETSPEAK, a tool

- □ to assist with word choice, and
- □ to check phrase commonness.

NETSPEAK implements wildcard queries on top of a Web n-gram index.

Netspeak The Writing Assistant

| looks fine ? me | Search |
|-----------------|--------|
|-----------------|--------|

| Frequency | | Phrase | Example |
|-----------|---------|---------------------------------|---------------|
| 19,103 | 93.8 % | looks fine <mark>to</mark> me | \pm |
| 810 | 4.0 % | looks fine <mark>for</mark> me | \pm |
| 353 | 1.7 % | looks fine <mark>with</mark> me | |
| 107 | 0.5 % | looks fine <mark>by</mark> me | Đ |
| 20,373 | 100.0 % | | 0.186 seconds |

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http://www.netspeak.cc

Given a set of *n*-grams, $n \leq 5$, and their frequencies.

A query q defines a pattern as a sequence of n-grams and wildcards.

A wildcard may be substituted for a defined subset of the *n*-grams.

Given a query q, retrieve all n-grams that match q.

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Given a query q, retrieve all n-grams that match q.

Straightforward solution:

- \Box Construct a keyword index for the *n*-grams.
- \Box Retrieve all *n*-grams that contain all of *q*'s words.
- \Box Compile a pattern matcher from q and filter the retrieved n-grams.

Improvements:

- \Box Exploit information encoded in queries and *n*-grams, and that *n* is small.
- \Box Exploit closed retrieval settings, e.g., the *n*-gram set is constant.
- □ Trade wildcard expressiveness and retrieval recall for time.
- Exploit information about the application domain.

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use the same ?

- □ Only 4-grams can match.
- **□** First word use, second word the, third word same.

Our index stores information about *n*-gram length and word position in the pre-image of the index lookup function.

prefer * over

- \square 2- to 5-grams can match.
- □ First word prefer, and last word over.

Variable-length queries are sub-divided into fixed-length queries: prefer over; prefer ? over; prefer ?? over; prefer ??? over

More search heuristics are described in [Stein *et al.*, ECIR'2010]

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Given a set G of n-grams, $n \leq 5$, and their frequencies $f : G \to \mathbf{N}$.

A query q defines a pattern as a sequence of n-grams and wildcards.

A wildcard may be substituted for every n-gram from a defined subset of G.

Given a query q, retrieve all n-grams R from G that match q.

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Given a query q, retrieve all n-grams R from G that match q.

Straightforward solution:

- \Box Construct an inverted index $\mu: V \to \mathcal{P}(G)$, where V is G's vocabulary.
- □ Retrieve all *n*-grams $R = \bigcap_{w \in q} \mu(w)$ that contain all of *q*'s words $w \in V$.
- \Box Compile a pattern matcher from q and filter R.

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NETSPEAK's approach:

 $\Box \text{ Construct an inverted index } \mu: V \times \underbrace{\{1, \dots, 5\}} \times \underbrace{\{1, \dots, 5\}} \to \mathcal{P}(G)$

□ Sort $\mu(w, i, j)$ in descending order of f, where $w \in V$ and $i, j \in \{1, ..., 5\}$.

n-gram length word position

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n-gram length word position

- □ Subdivide *q* into $\{q_1, \ldots, q_m\}$ so that $R_q = \bigcup_{i=1}^m R_{q_i}$, and each q_i matches only *n*-grams with a fixed length. Process the sub-queries in parallel.
- □ Retrieve all *n*-grams $R_{q_i} = \bigcap_{w \in q_i} \mu(w, |q_i|, q_{i|w})$, where $q_{i|w}$ is *w*'s position in q_i .

NETSPEAK's approach:

- $\Box \text{ Construct an inverted index } \mu: V \times \underbrace{\{1, \dots, 5\}}_{\text{n-gram length}} \times \underbrace{\{1, \dots, 5\}}_{\text{word position}} \to \mathcal{P}(G)$
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- □ Start to process each $\mu(w, i, j)$ at entry k, with $f(\mu(w, i, j)_k) \le \min_{g \in q} (f(g))$.
- $\hfill\square$ Stop to process each $\mu(w,i,j)$ at entry

 $\left\{ \begin{array}{ll} |\mu(w,i,j)| & \text{if the postlist is smaller than a page, or} \\ l_1 & \text{if a pre-specified amount of results have been retrieved, or} \\ l_2 & \text{if } \sum_{i'=0}^{l_2} f(\mu(w,i,j)_{i'}) \text{ covers } \kappa\% \text{ of the frequency distribution.} \end{array} \right.$